



## Introduction

- **QSVT**: Unification of quantum algorithms.
- **Challenge**: Ultimately relies on a block encoding circuit for the system matrix.
  - > Not immediately clear how to construct these.
  - > See: talk by [Chao Yang](#) on Thursday for strategies to construct  $\mathcal{O}(\text{polylog}(N))$  circuits.

## FABLE

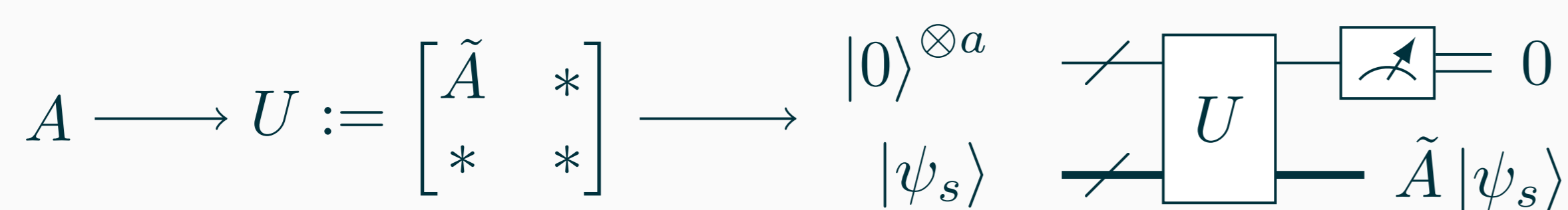
FABLE circuits ...

- generically contain:
  - >  $N^2$  two-qubit CNOTs and  $N^2$  single-qubit  $R_y$  gates.
  - >  $\mathcal{O}(\text{polylog}(N))$  overhead.
- are **compressible**, in practice often more efficient **approximate circuits**.
- are **easy** to generate and simulate in **QCLAB**:

```
1 [[circuit, ~, alpha] = fable(A);
2 U = circuit.matrix;
3 A_BE = 2^n * U(1:2^n, 1:2^n);
4 norm(A - alpha*A_BE)      ans =
                               2.7439e-15
```

## Block encodings

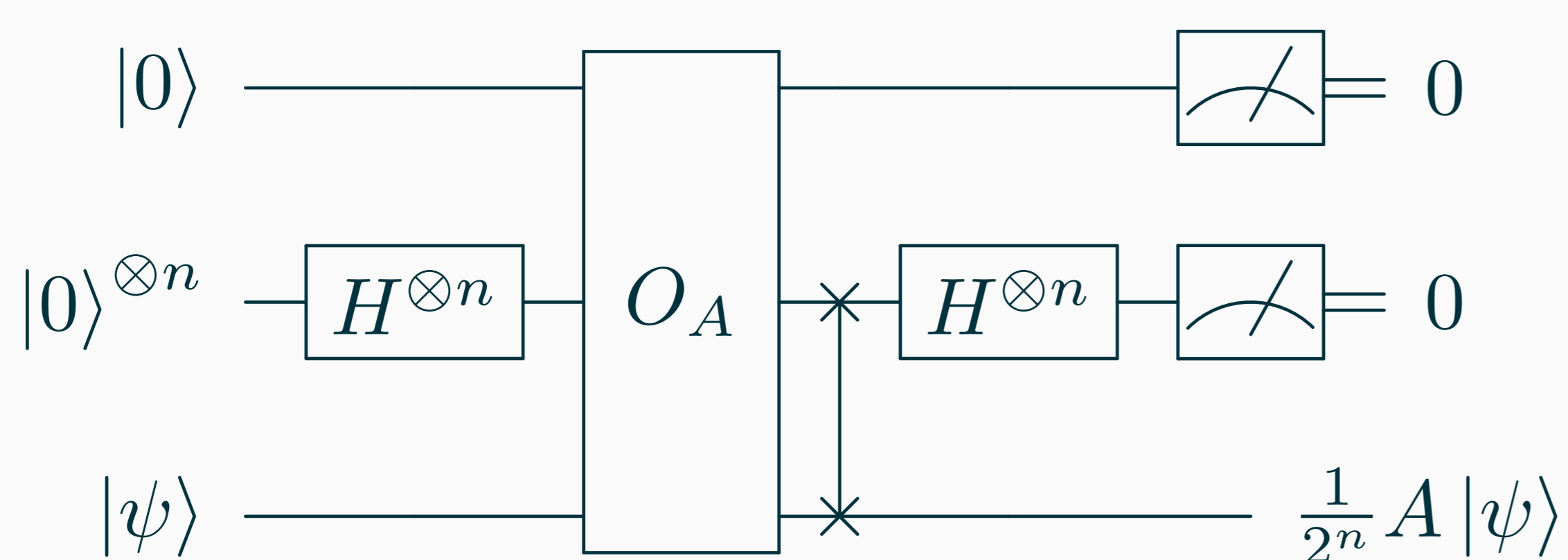
A **block encoding** approximately embeds a –not necessarily unitary– operator  $A$  as the leading principal block in a larger unitary  $U$ :



with  $U$  an  $(\alpha, a, \epsilon)$ -block encoding of  $A$  where  $\|A - \alpha \tilde{A}\|_2 \leq \epsilon$ .

## Circuit template

- FABLE circuits use the following **template**:



- $O_A$  is a **matrix query oracle**:

$$O_A |0\rangle |i\rangle |j\rangle = (a_{ij} |0\rangle + \sqrt{1 - |a_{ij}|^2} |1\rangle) |i\rangle |j\rangle$$

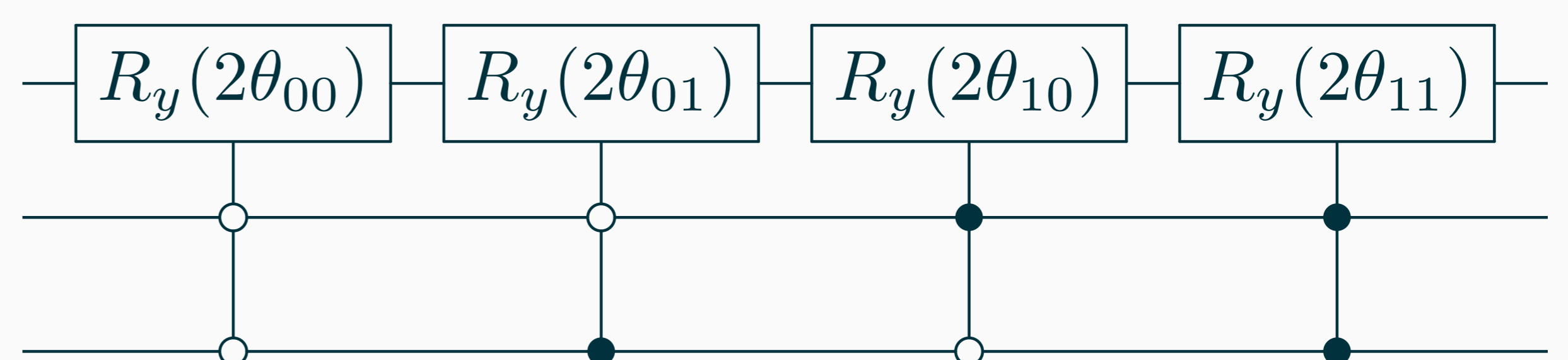
- FABLE solves the problem of **synthesizing  $O_A$  in simple one- and two-qubit gates for arbitrary matrices  $A$** .
- Hadamard and SWAP gates add  $\mathcal{O}(\text{polylog}(N))$  overhead.
- The result is an  $(1/2^n, n + 1, 0)$ -block-encoding for an  $n$ -qubit matrix  $A$  with  $|a_{ij}| \leq 1$ .

## $O_A$ for real-valued data

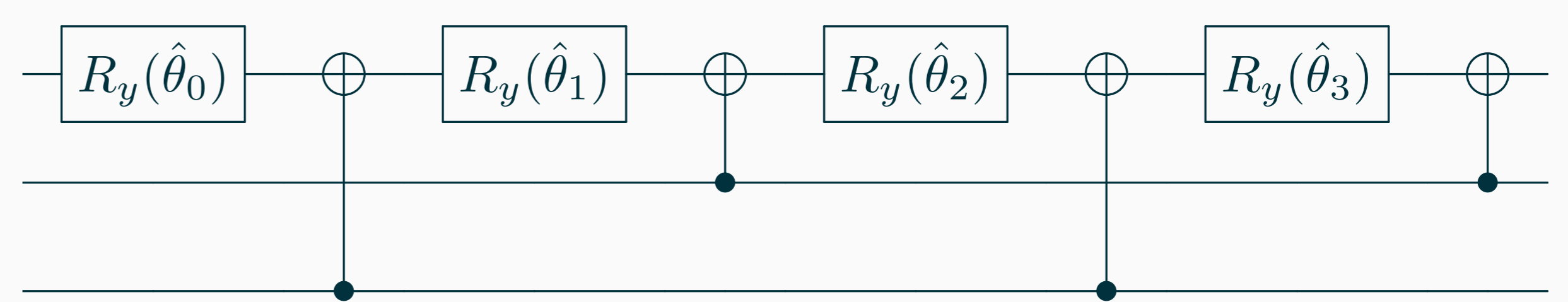
- For an entry  $a_{ij}$ ,  $O_A$  acts as an  $R_y$  gate with angle  $\theta_{ij} = \arccos(a_{ij})$ :

$$R_y(2\theta_{ij}) |0\rangle = \begin{bmatrix} \cos(\theta_{ij}) & -\sin(\theta_{ij}) \\ \sin(\theta_{ij}) & \cos(\theta_{ij}) \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} a_{ij} \\ \sqrt{1 - a_{ij}^2} \end{bmatrix}$$

- Naive implementation of  $O_A$  requires  $\mathcal{O}(N^4)$  elementary gates:



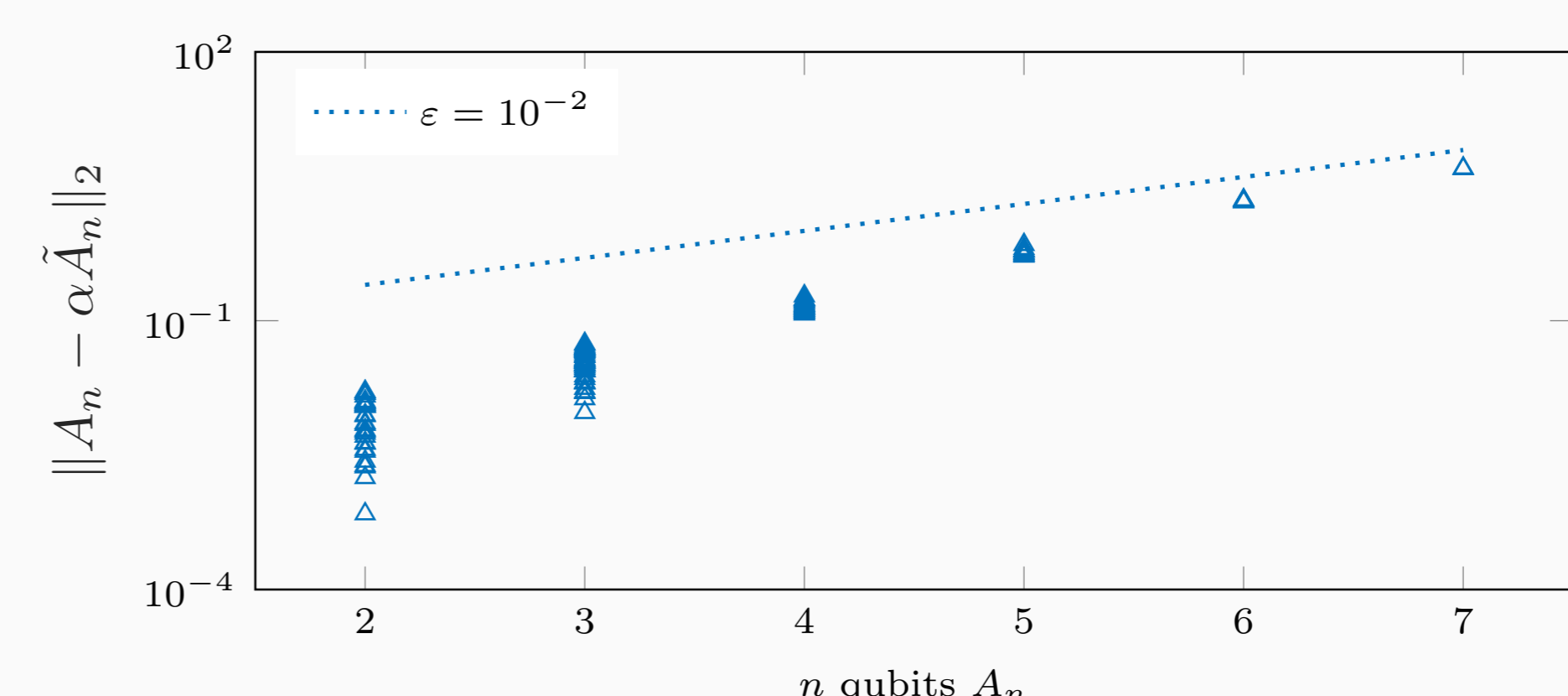
- An improved implementation reduces cost to  $\mathcal{O}(N^2)$ :



- The angles  $\theta$  and  $\hat{\theta}$  are related by  $(\hat{H}^{\otimes 2n} P_G) \hat{\theta} = \theta$ , which can be solved classically with cost  $\mathcal{O}(N^2 \log N^2)$ .

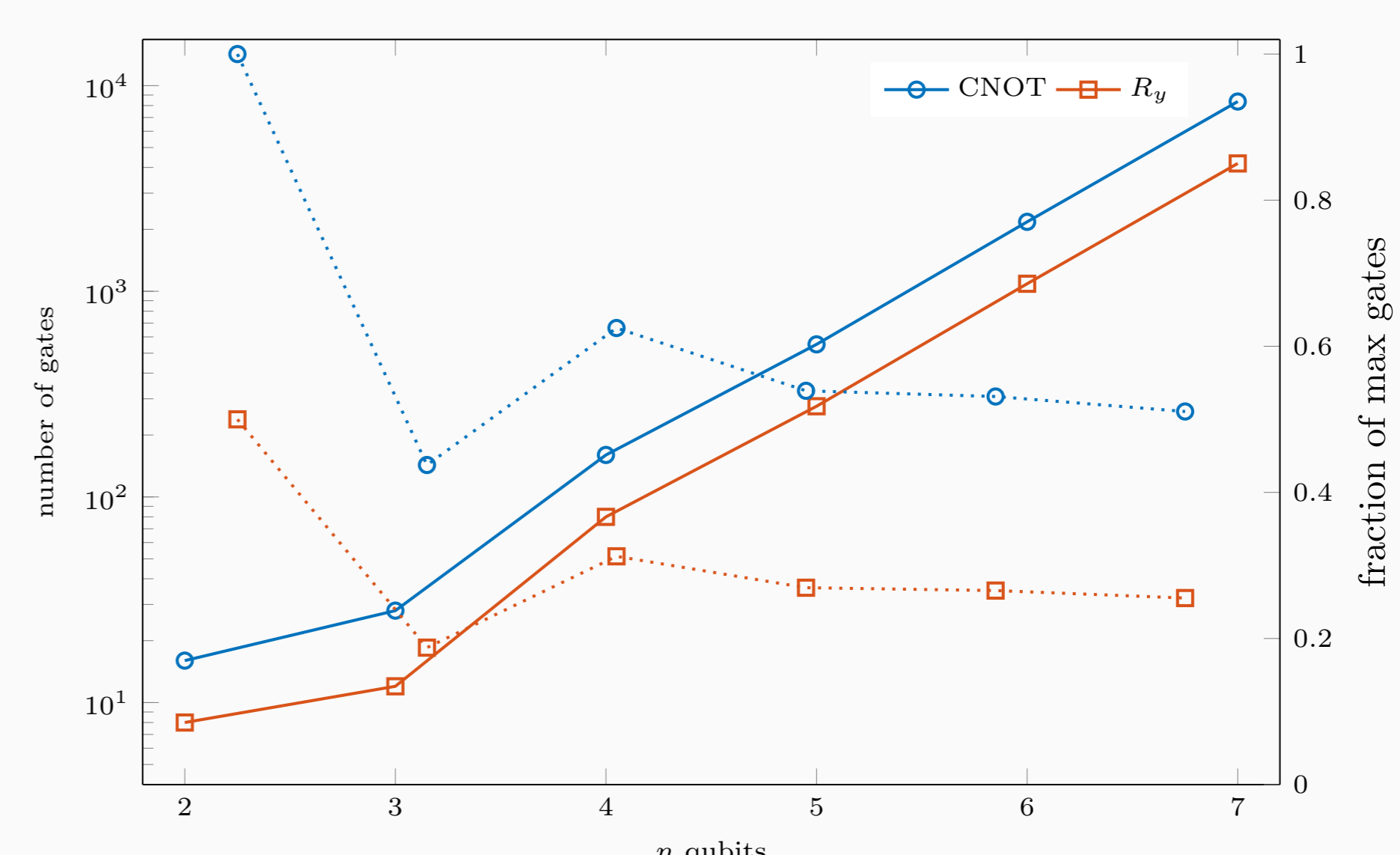
## Compression Theorem

For a given cutoff threshold  $\epsilon \geq 0$ , setting all angles  $|\hat{\theta}_i| \leq \epsilon$  to zero, results in a  $(1/2^n, n + 1, 2\pi N\epsilon)$ -block-encoding of an  $n$ -qubit matrix  $A$ .



## Heisenberg XXX Hamiltonian

The Heisenberg XXX model can be encoded w/o compression ( $\epsilon = 0$ ), with 50% fewer CNOTs and 75% fewer  $R_y$  gates:



## Conclusion

- FABLE can easily generate and compress block encoding circuits up to  $\sim 15$  qubit operators.
- Encoding of complex-valued data doubles the number of gates.

**Reference:** Camps D. and Van Beeumen R., *Fast Approximate Quantum Circuit Synthesis for Block Encoding Matrices*, In Preparation.